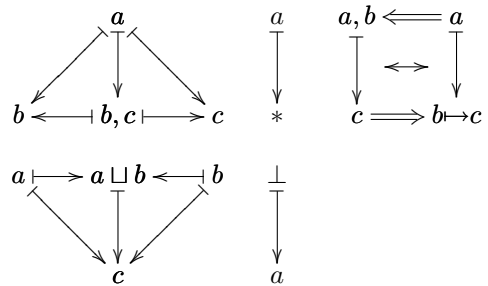


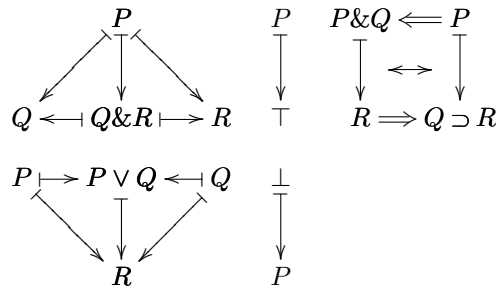
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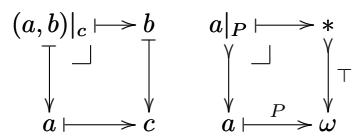
The BiCCC structure and the classifier axiom
CCC:



HA:



Topos:



The pre-hyperdoctrine structure

Hyperdoctrine:

$$\begin{array}{ccccc}
 & \begin{array}{c} (P) \\ (a, b) \end{array} & \xrightarrow{\text{co}\square} & \begin{array}{c} (b=b' \& P) \\ (a, b, b') \end{array} & \begin{array}{c} (P) \\ (a, b) \end{array} & \xrightarrow{\text{co}\square} & \begin{array}{c} (\exists b.P) \\ (a) \end{array} \\
 & \downarrow & & \downarrow & \downarrow & & \downarrow \\
 \begin{array}{c} (P) \\ (a) \end{array} & \xrightarrow{\square} & \begin{array}{c} (P) \\ (b) \end{array} & \begin{array}{c} (Q[b, b]) \\ (a, b) \end{array} & \xrightarrow{\square} & \begin{array}{c} (Q[b, b']) \\ (a, b, b') \end{array} & \begin{array}{c} (Q) \\ (a, b) \end{array} & \xrightarrow{\square} & \begin{array}{c} (Q) \\ (a) \end{array} \\
 & & & \downarrow & & \downarrow & \downarrow & & \downarrow \\
 & & & \begin{array}{c} (R) \\ (a, b) \end{array} & & & \begin{array}{c} (R) \\ (a, b) \end{array} & & \begin{array}{c} (\forall b.R) \\ (a) \end{array} \\
 \\
 a \vdash \longrightarrow b & & a, b \vdash \longrightarrow a, b, b' & & a, b \vdash \longrightarrow a
 \end{array}$$

LCCC:

$$\begin{array}{ccccc}
 & \begin{array}{c} (c) \\ (a, b) \end{array} & \xrightarrow{\text{co}\square} & \begin{array}{c} ((b=b'), c) \\ (a, b, b') \end{array} & \begin{array}{c} (c) \\ (a, b) \end{array} & \xrightarrow{\text{co}\square} & \begin{array}{c} (b, c) \\ (a) \end{array} \\
 & \downarrow & & \downarrow & \downarrow & & \downarrow \\
 \begin{array}{c} (c) \\ (a) \end{array} & \xrightarrow{\square} & \begin{array}{c} (c) \\ (b) \end{array} & \begin{array}{c} (d) \\ (a, b) \end{array} & \xrightarrow{\square} & \begin{array}{c} (d) \\ (a, b, b') \end{array} & \begin{array}{c} (d) \\ (a, b) \end{array} & \xrightarrow{\square} & \begin{array}{c} (d) \\ (a) \end{array} \\
 & & & \downarrow & & \downarrow & \downarrow & & \downarrow \\
 & & & \begin{array}{c} (e) \\ (a, b) \end{array} & & & \begin{array}{c} (e) \\ (a, b) \end{array} & & \begin{array}{c} (b \dashv\vdash e) \\ (a) \end{array} \\
 \\
 a \vdash \longrightarrow b & & a, b \vdash \longrightarrow a, b, b' & & a, b \vdash \longrightarrow a
 \end{array}$$

Two CCompC structures in a topos

The two CCompC structures in a topos:

$$\begin{array}{ccc}
 \begin{array}{c} (P) \\ (a) \end{array} \multimap \begin{array}{c} (\top) \\ (b) \end{array} \multimap \begin{array}{c} (Q) \\ (c) \end{array} & & \begin{array}{c} (b) \\ (a) \end{array} \multimap \begin{array}{c} (*) \\ (c) \end{array} \multimap \begin{array}{c} (e) \\ (d) \end{array} \\
 \Downarrow \quad \Updownarrow \quad \Uparrow \quad \Downarrow & & \Downarrow \quad \Updownarrow \quad \Uparrow \quad \Downarrow \\
 a \multimap b \multimap c|_Q & & a \multimap c \multimap d, e
 \end{array}$$

Cartesian morphisms project into pullbacks:

$$\begin{array}{ccc}
 \begin{array}{c} (P) \\ (a) \end{array} & & \begin{array}{c} (c) \\ (a) \end{array} \\
 \swarrow \quad \Downarrow \quad \searrow & & \swarrow \quad \Downarrow \quad \searrow \\
 a|_P \multimap a & & a, c \multimap a \\
 \swarrow \quad \Downarrow \quad \searrow & & \swarrow \quad \Downarrow \quad \searrow \\
 b|_P \multimap b & & b, c \multimap b
 \end{array}$$

Cocartesian morphisms induce isos and epis:

$$\begin{array}{ccc}
 \begin{array}{c} (P) \\ (a, b) \end{array} & & \begin{array}{c} (P) \\ (a, b) \end{array} \\
 \swarrow \quad \Downarrow \quad \searrow \text{co}\square & & \swarrow \quad \Downarrow \quad \searrow \text{co}\square \\
 a, b|_P \multimap a, b & & a, b|_P \multimap a, b \\
 \swarrow \quad \Downarrow \quad \searrow & & \swarrow \quad \Downarrow \quad \searrow \\
 a|\exists b.P \multimap a & & a, b, b'|_{b=b' \& P} \multimap a, b, b' \\
 \begin{array}{c} (c) \\ (a, b) \end{array} & & \begin{array}{c} (c) \\ (a, b) \end{array} \\
 \swarrow \quad \Downarrow \quad \searrow \text{co}\square & & \swarrow \quad \Downarrow \quad \searrow \text{co}\square \\
 (a, b), c \multimap a, b & & a, b, c \multimap a, b \\
 \swarrow \quad \Downarrow \quad \searrow & & \swarrow \quad \Downarrow \quad \searrow \\
 a, (b, c) \multimap a & & a, b, b', (b=b'), c \multimap a, b, b' \\
 \begin{array}{c} (\exists b.P) \\ (a) \end{array} & & \begin{array}{c} (b=b' \& P) \\ (a) \end{array} \\
 \Downarrow & & \Downarrow \\
 a & & a
 \end{array}$$

Basic constructions with NNOs

$$\begin{array}{c}
 * \xrightarrow{\quad} 0 \xrightarrow{\quad} 1 \xrightarrow{\quad} 2 \\
 \searrow \quad \downarrow \quad \downarrow \quad \downarrow \\
 (n \mapsto n) \mapsto (n \mapsto n+1) \mapsto (n \mapsto n+2)
 \end{array}
 \qquad
 \begin{array}{c}
 m \\
 \downarrow \\
 n \mapsto n+m
 \end{array}$$

$$\begin{array}{c}
 * \xrightarrow{\quad} 0 \xrightarrow{\quad} 1 \xrightarrow{\quad} 2 \\
 \searrow \quad \downarrow \quad \downarrow \quad \downarrow \\
 (n \mapsto 0) \mapsto (n \mapsto n) \mapsto (n \mapsto 2n)
 \end{array}
 \qquad
 \begin{array}{c}
 m \\
 \downarrow \\
 n \mapsto mn
 \end{array}$$

$$\begin{array}{c}
 * \xrightarrow{\quad} 0 \xrightarrow{\quad} 1 \xrightarrow{\quad} 2 \\
 \searrow \quad \downarrow \quad \downarrow \quad \downarrow \\
 (n \mapsto 1) \mapsto (n \mapsto n) \mapsto (n \mapsto n^2)
 \end{array}
 \qquad
 \begin{array}{c}
 m \\
 \downarrow \\
 n \mapsto n^m
 \end{array}$$

$$\begin{array}{c}
 * \xrightarrow{\quad} 0 \xrightarrow{\quad} 1 \xrightarrow{\quad} 2 \\
 \searrow \quad \downarrow \quad \downarrow \quad \downarrow \\
 (n \mapsto n) \mapsto (n \mapsto n-1) \mapsto (n \mapsto n-2)
 \end{array}
 \qquad
 \begin{array}{c}
 m \\
 \downarrow \\
 n \mapsto n-m
 \end{array}$$

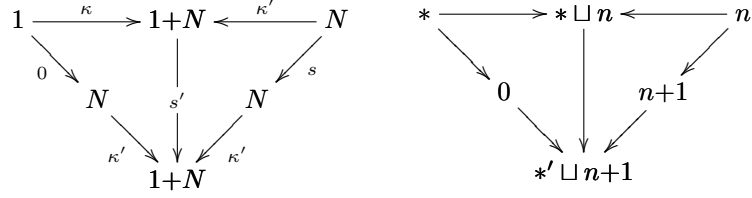
NNOs: the morphism p'

Fact: in a topos with NNO the map $[0, s] : 1+N \rightarrow N$ is an iso.

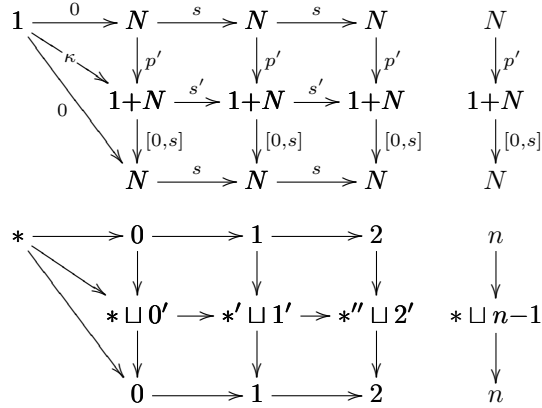
First we need to define the arrow $s' : 1+N \rightarrow 1+N$,
using a factorization through a coproduct.

Note that s' takes the ' $*$ ' in ' $* \sqcup n$ ' to 0, not to $*$ '.

$s' := [0; \kappa', s; \kappa']$.



Now we can define $p' : N \rightarrow 1+N$ by factoring (κ, s') through the NNO.
It is possible to show that p' and $[0, s]$ are inverses.



Define $n \mapsto n-1$ as $p'; [0, \text{id}]$,

The arrow $m \mapsto (n - m)$ of the previous page